Problem 1. The snapshot does not need to be atomic. Even if a non-atomic snapshot is used in the scan routine (i.e., the process epochs were never simultaneously equal to their values in the node epoch vector), this cannot cause a node to be incorrectly reclaimed (freed when some process still holds references to it).

Problem 2. See the original Hazard Pointers paper:


In particular, see the Scan algorithm in Figure 3. Essentially, the idea is to add all hazard pointers in the system to a hash map (which takes $O(n)$ time) and then check for every node in the limbo list if it is protected by searching for that node in the hash map ($m$ queries that take $O(1)$ time each).