

The Power of Registers

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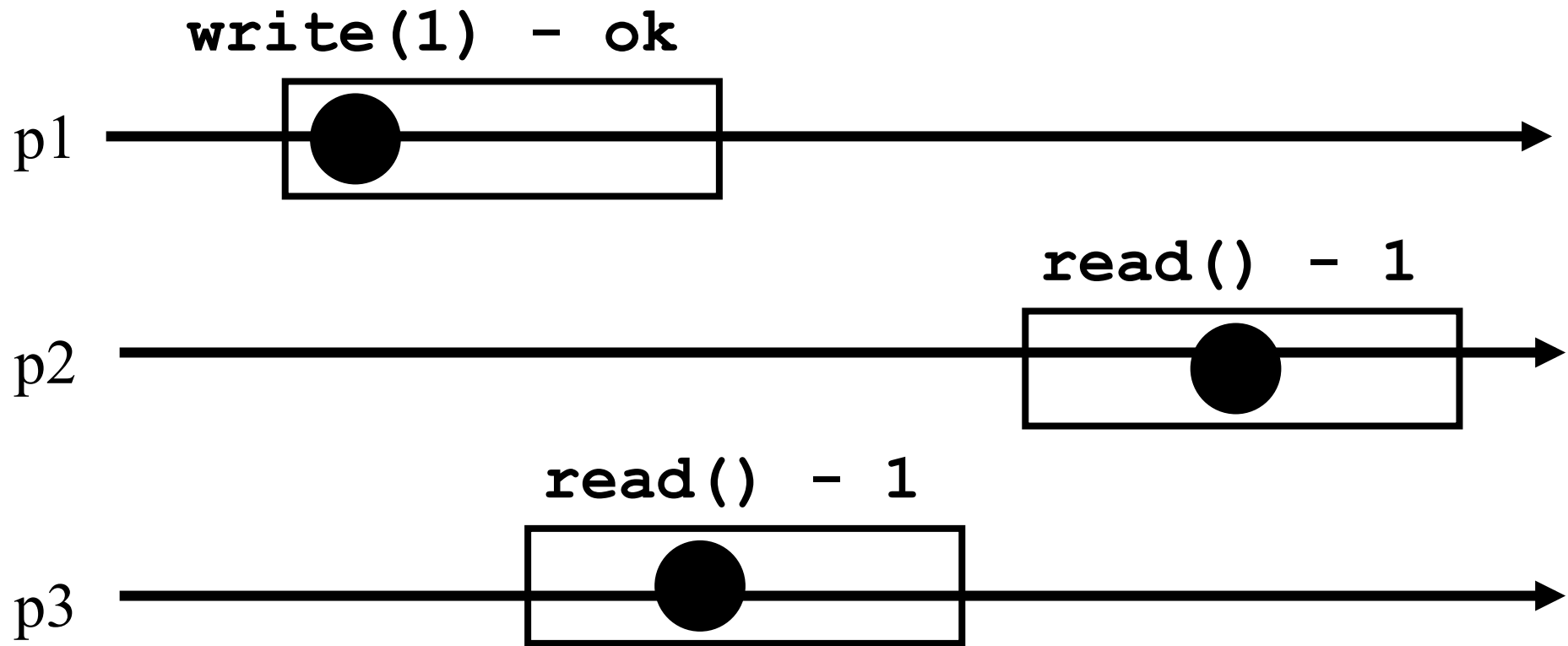


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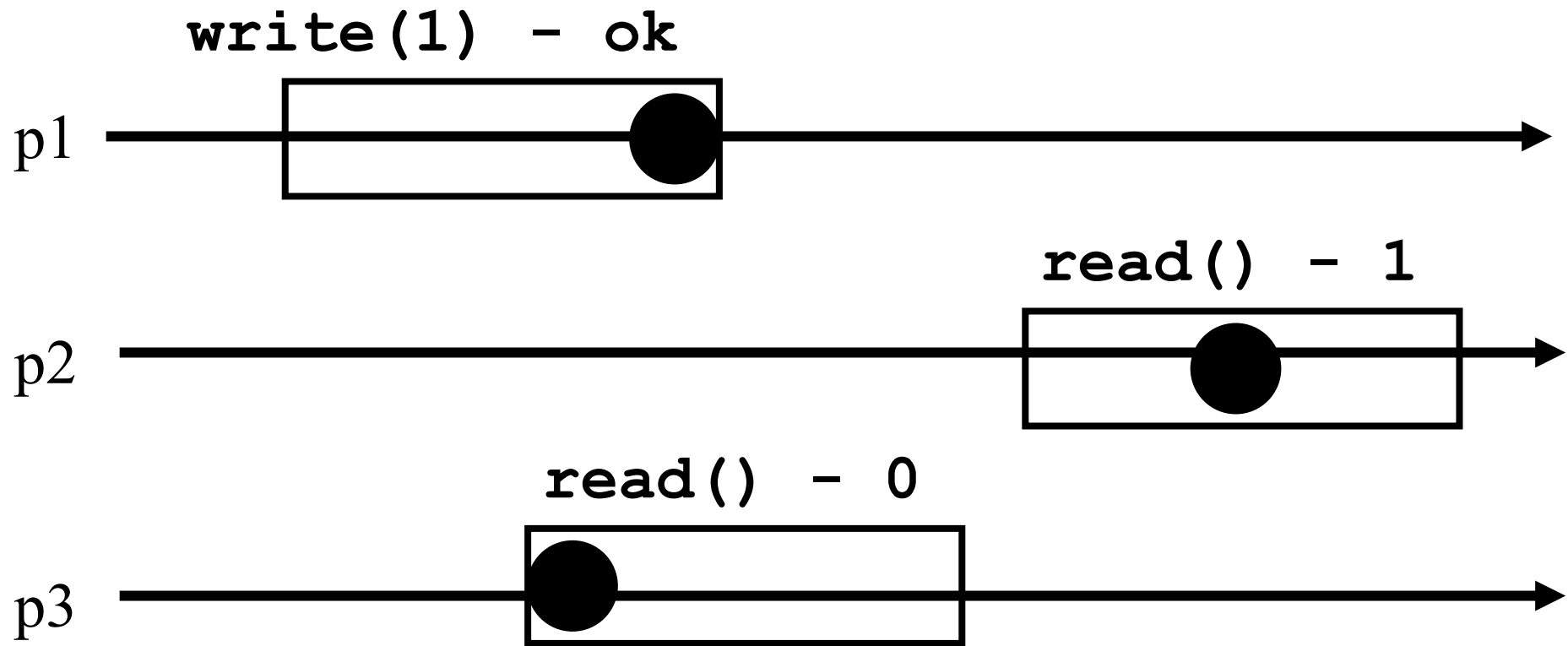
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Atomic execution



Atomic execution



Registers

- **Question 1:** what objects can we implement with registers?
- Question 2: what objects we cannot implement?

Wait-free implementations of atomic objects

- An **atomic** object is simply defined by its sequential specification; i.e., by how its operations should be implemented when there is no concurrency
- Implementations should be **wait-free**: every process that invokes an operation eventually gets a reply (unless the process crashes)

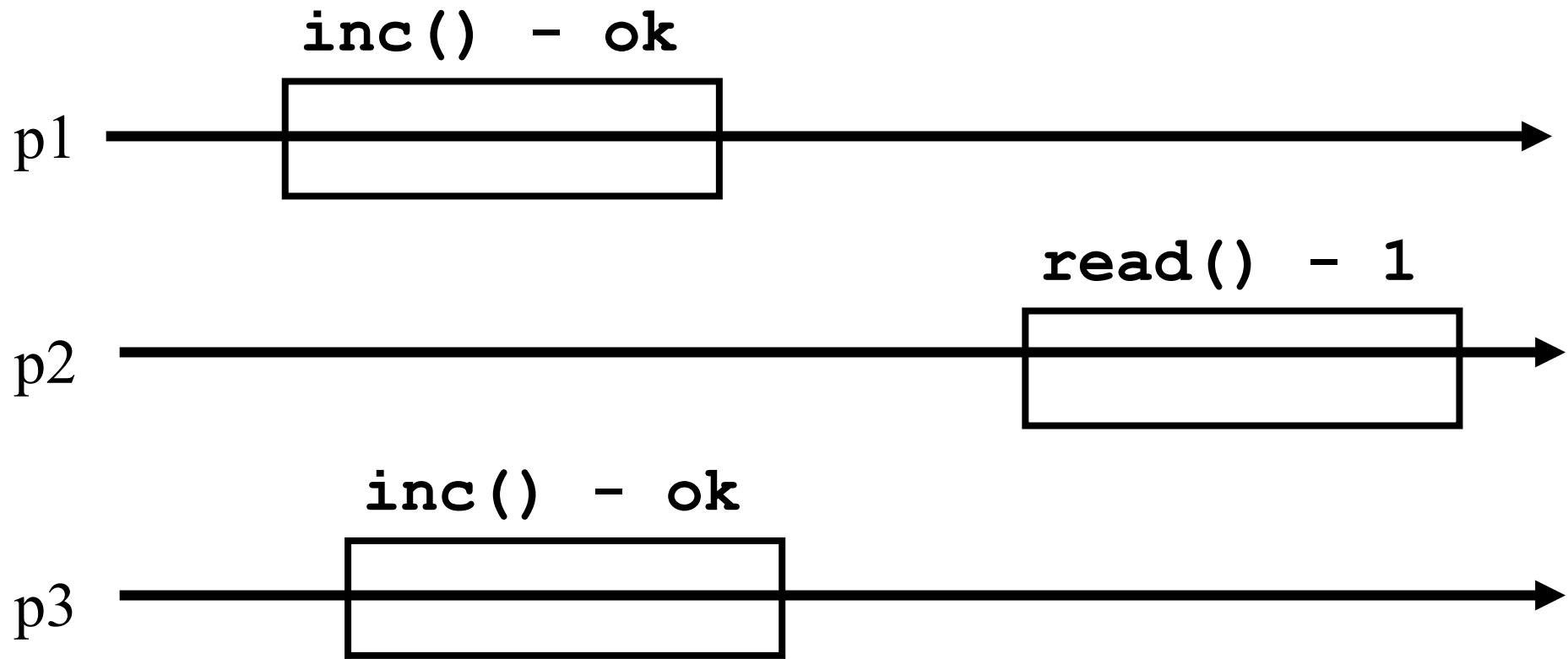
Counter (sequential spec)

- A **counter** has two operations ***inc()*** and ***read()*** and maintains an integer *x* *init to 0*
- ***read():***
 - return(x)
- ***inc():***
 - $x := x + 1;$
 - return(ok)

Naive implementation

- The processes share one register Reg
- ***read()***:
 - return(Reg.read())
- ***inc()***:
 - temp:= Reg.read()+1;
 - Reg.write(temp);
 - return(ok)

Atomic execution?



Atomic implementation

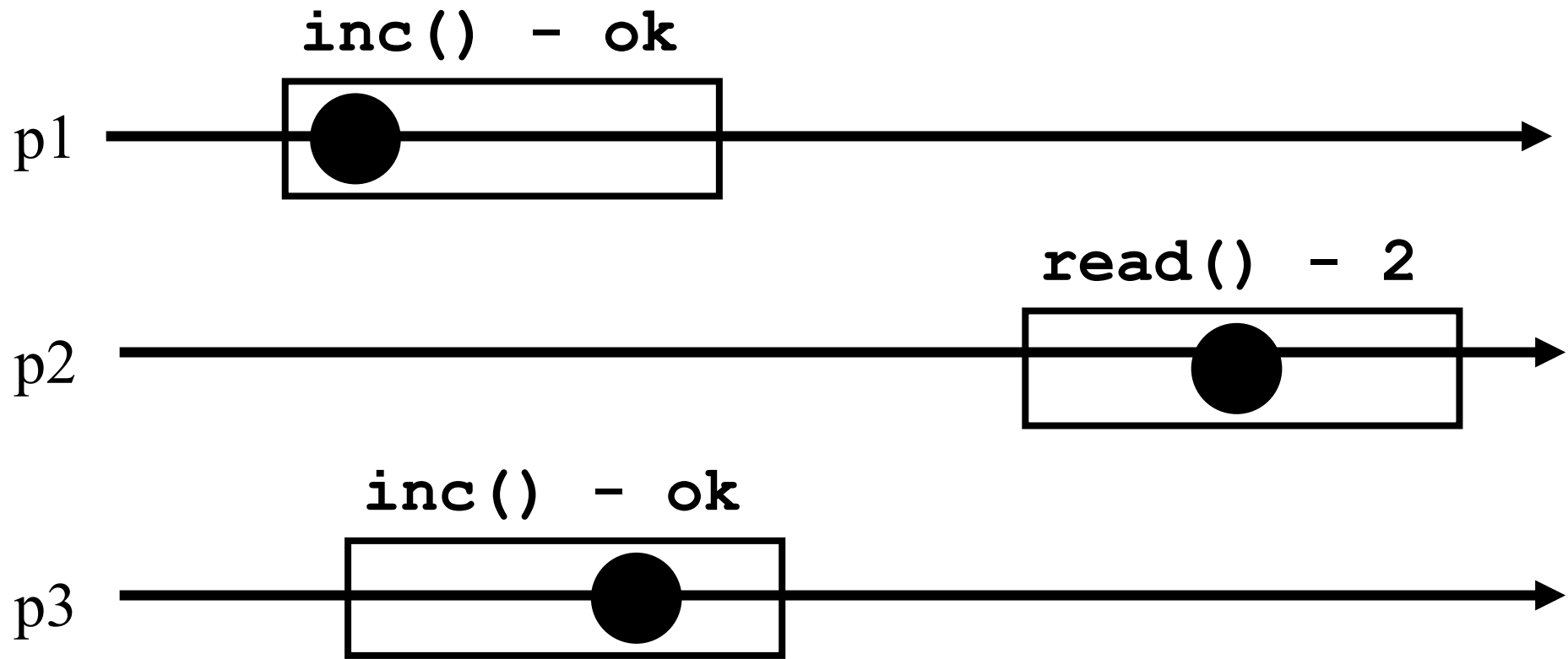
- The processes share an array of registers $\text{Reg}[1, \dots, n]$
- ***inc()***:
 - $\text{Reg}[i].\text{write}(\text{Reg}[i].\text{read()} + 1);$
 - $\text{return}(\text{ok})$

Atomic implementation

• *read():*

- sum := 0;
- for j = 1 to n do
 - sum := sum + Reg[j].read();
- return(sum)

Atomic execution?



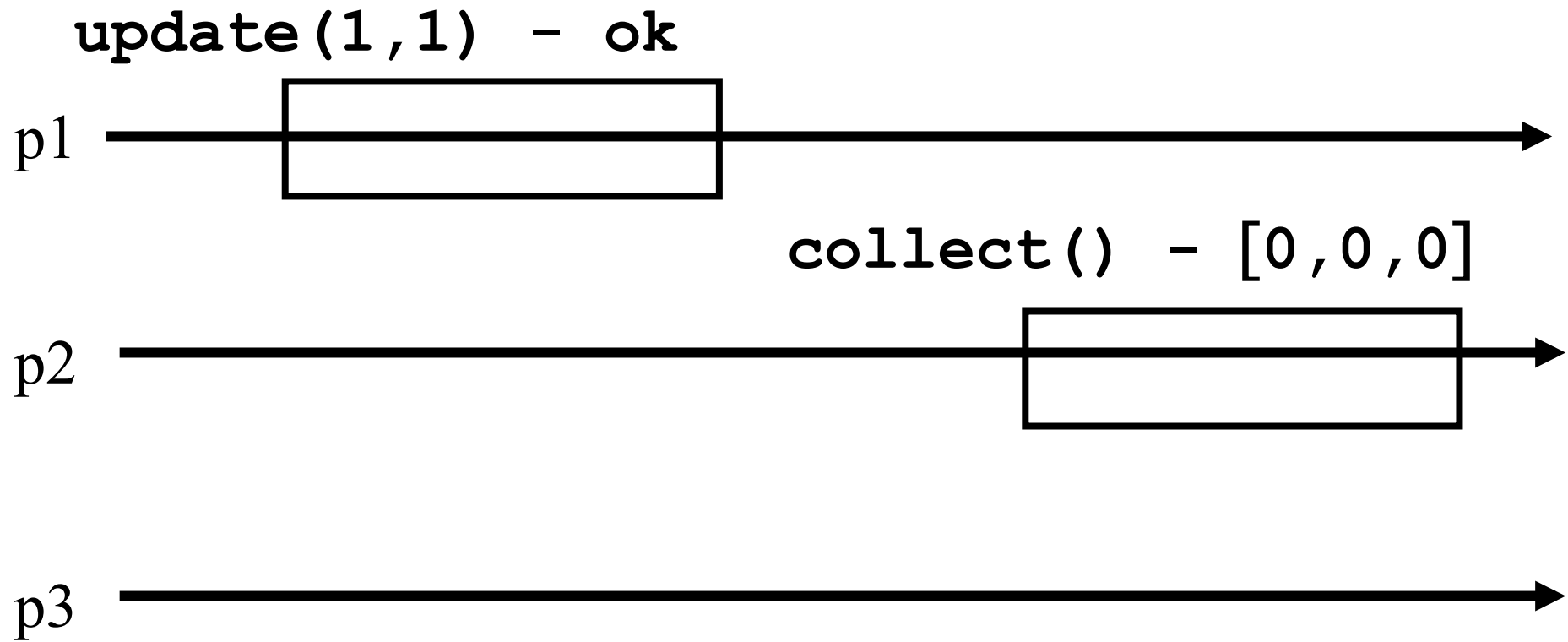
Snapshot (sequential spec)

- A **snapshot** has operations **update()** and **scan()** and maintains an array x of size n
- **scan():**
 - return(x)
- **update(i, v):**
 - $x[i] := v;$
 - return(ok)

Very naive implementation

- Each process maintains an array of integer variables x init to $[0, \dots, 0]$
- ***scan()***:
 - return(x)
- ***update(i,v)***:
 - $x[i] := v;$
 - return(ok)

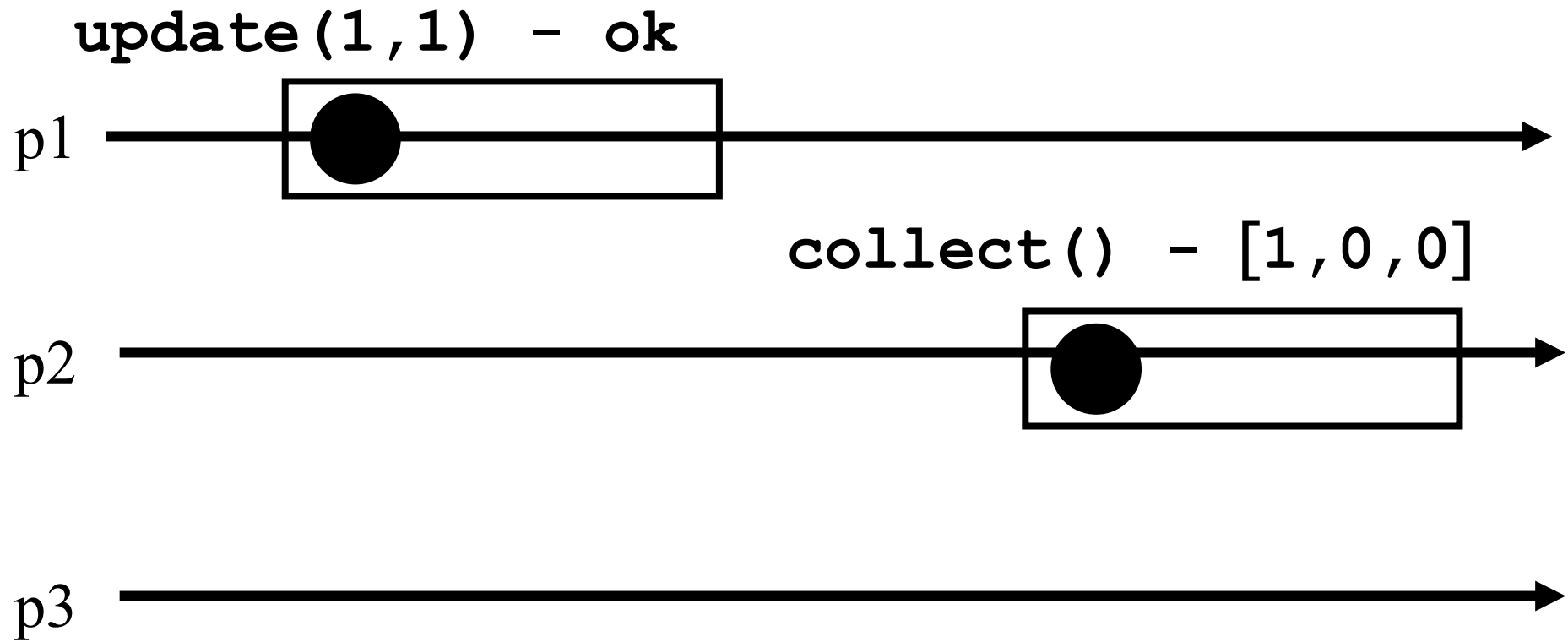
Atomic execution?



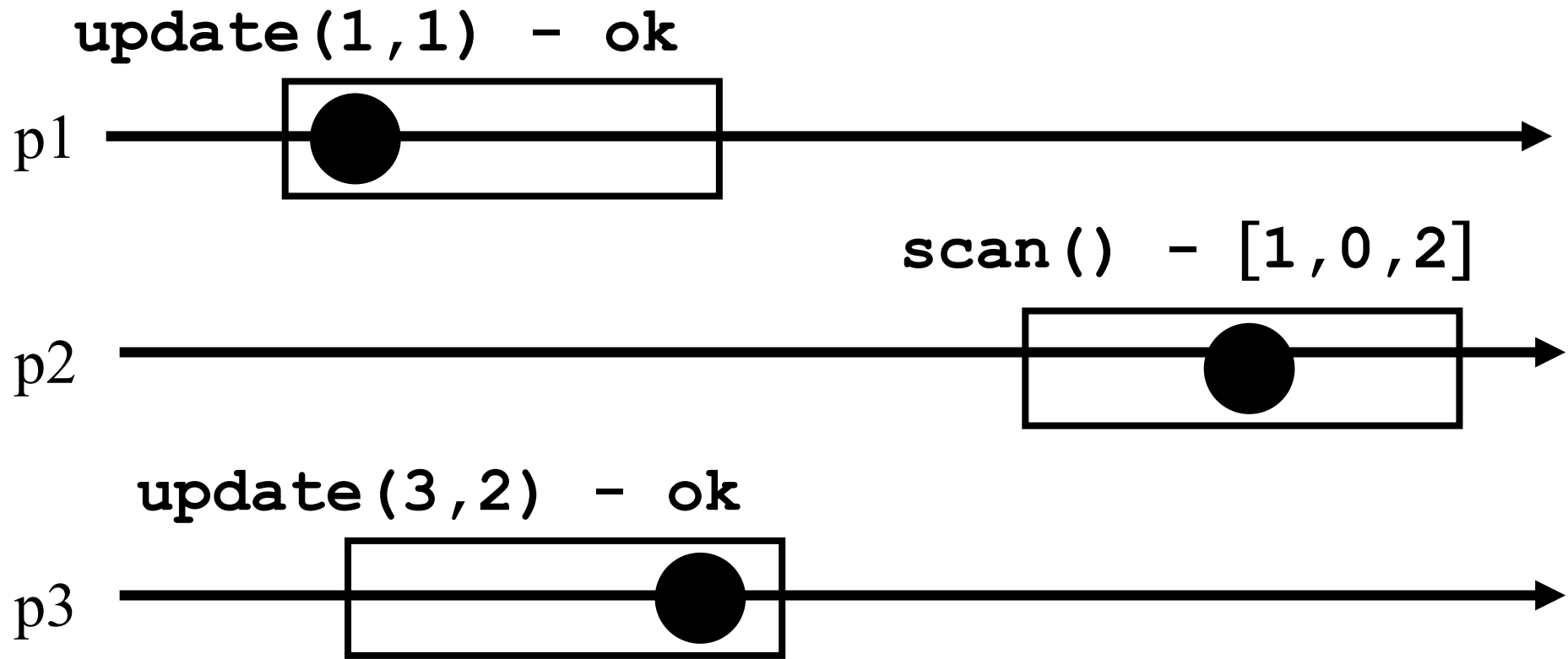
Less naive implementation

- The processes share one array of N registers
Reg[1,..,N]
- **scan():**
 - for j = 1 to N do
 - x[j] := Reg[j].read();
 - return(x)
- **update(i,v):**
 - Reg[i].write(v); return(ok)

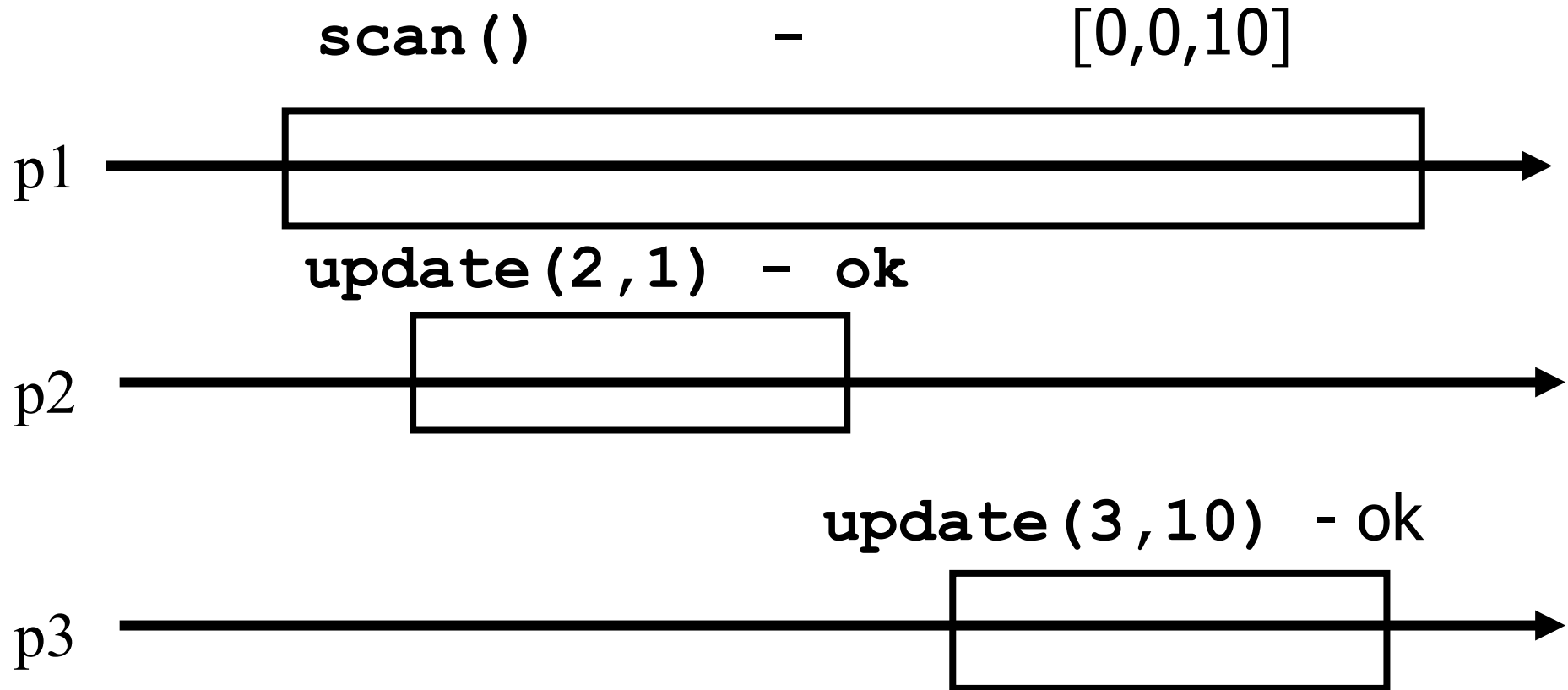
Atomic execution?



Atomic execution?



Atomic execution?



Non-atomic vs atomic snapshot

- What we implement here is some kind of **regular** snapshot:
 - A **scan** returns, for every index of the snapshot, the last written values or the value of any concurrent update
 - We call it **collect**

Key idea for atomicity

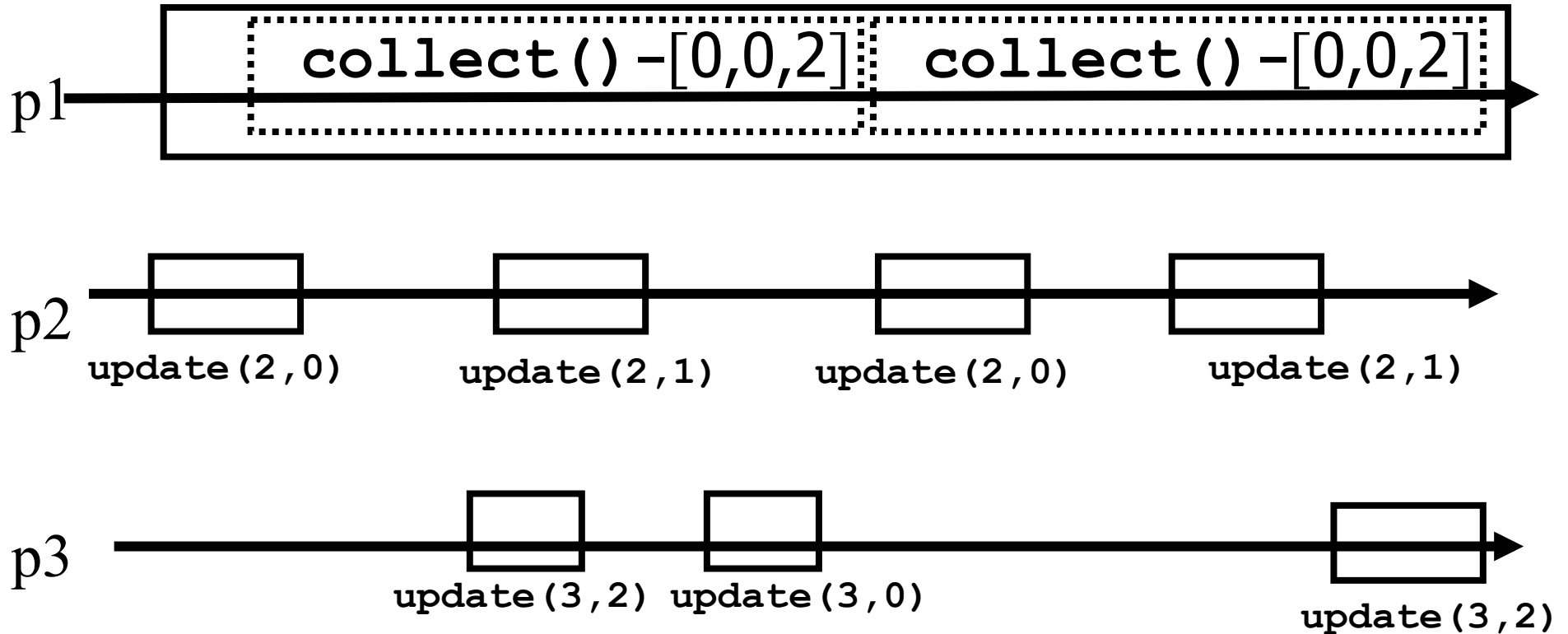
- To ***scan***, a process keeps reading the entire snapshot (i.e., it ***collect***), until two results are the **same**
- This means that the snapshot did not change, and it is safe to return without violating atomicity

Same value vs. Same timestamp

scan ()

-

[0,0,2]



Enforcing atomicity

- The processes share one array of N registers $\text{Reg}[1,\dots,N]$; each contains a value and a timestamp
- We use the following operation for modularity
- ***collect()***:
 - for $j = 1$ to N do
 - $x[j] := \text{Reg}[j].\text{read}()$;
 - return(x)

Enforcing atomicity (cont'd)

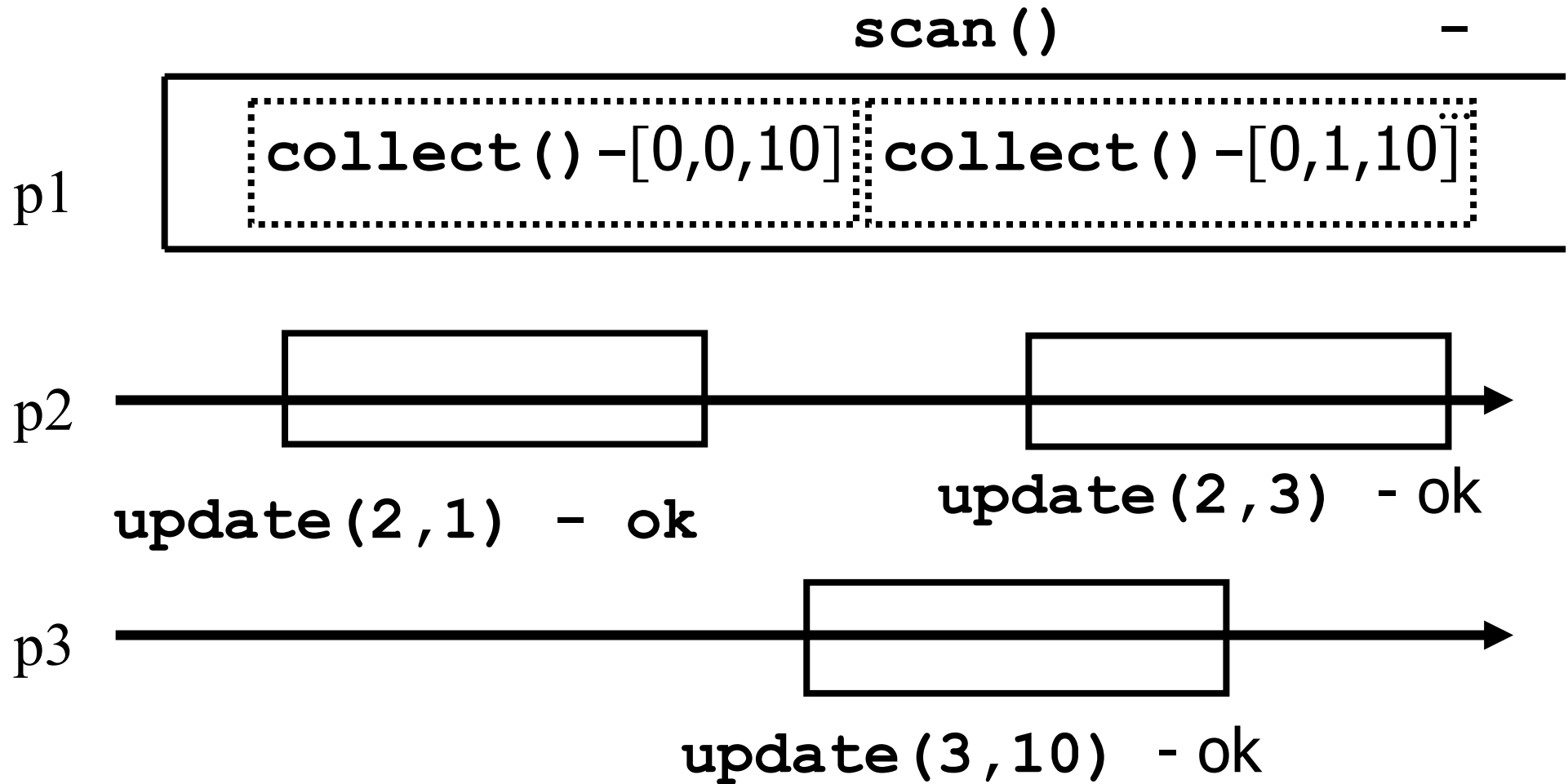
☛ ***scan()***:

- ☛ temp1 := self.collect();
- ☛ while(true) do
 - ☛ temp2 := self.collect();
 - ☛ if (temp1 = temp2) then
 - ☛ return (temp1.val)
 - ☛ temp1 := temp2;

☛ ***update(i,v)***:

- ☛ ts := ts + 1;
- ☛ Reg[i].write(v,ts);
- ☛ return(ok)

Wait-freedom?



Key idea for atomicity & wait-freedom

- The processes share an array of ***registers*** $\text{Reg}[1,\dots,N]$ that contains each:
 - a value,
 - a timestamp, and
 - a copy of the entire array of values

Key idea for atomicity & wait-freedom (cont'd)

- To ***scan***, a process keeps collecting and returns a collect if it did not change, or some collect returned by a concurrent ***scan***
 - Timestamps are used to check if the collect changes or if a scan has been taken in the meantime
- To ***update***, a process ***scans*** and writes the value, the new timestamp and the result of the scan

Snapshot implementation

Every process keeps a local timestamp ts

• ***update(i, v):***

- $ts := ts + 1;$
- $Reg[i].write(v, ts, self.scan());$
- $return(ok)$

Snapshot implementation

• *scan()*:

- `t1 := self.collect(); t2 := t1`
- `while(true) do`
 - `t3 := self.collect();`
 - `if (t3 = t2) then return (t3);`
 - `for j = 1 to N do`
 - `if(t3[j,2] ≥ t1[j,2]+2) then`
 - `return (t3[j,3])`
 - `t2 := t3`

**Return the
first value in
each cell in t3**

Possible execution?

