Concurrent Algorithms

November 18, 2019

Exercise 8

(Given as an exam problem in 2016-2017)

Problem 1. Consider the following implementation of an obstruction-free consensus object from atomic multi-valued MRMW shared registers in a system of *n* processes. A process's id is known to itself as *i*.

```
Using: an array of atomic multi-valued MRMW shared registers T[1, 2, ..., n],
       initialized to 0;
Using: an array of atomic multi-valued MRMW shared registers V[1, 2, ..., n],
       initialized to (\bot, 0);
propose(v) {
    ts := i;
    while (true) do{
        T[i].write(ts);
        maxts := 0;
        val := \bot;
        for j = 1 to n do
            (t, vt) := V[j].read();
            if maxts < t then
                maxts := t;
                val := vt;
         if val = \bot then val := v;
         maxts := 0;
         for j = 1 to n do
             t := T[j].read();
             if maxts < t then maxts := t;
         if ts = maxts then
             V[i].write(val, ts);
             return(val);
         ts := ts + n;
    }
}
```

Recall that obstruction-free consensus ensures the property of *obstruction-freedom* instead of *wait-freedom*. **Your tasks:**

- 1. Explain what is obstruction-freedom and what is the difference between obstruction-freedom and wait-freedom.
- 2. Answer whether the implementation satisfies obstruction-freedom. Justify your answer.
- 3. Is the algorithm correct? If the algorithm is correct prove its correctness. Otherwise, provide an execution that shows the algorithm is incorrect.