

## Exercise 11

**Problem 1.** Consider the Disk Paxos algorithm in slides 14-15 of the lecture. The algorithm is reproduced below. If we omit line 11, is the algorithm still correct? Why or why not?

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**Algorithm 1** Obstruction-free consensus with Memory Failures
 

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1: procedure PROPOSE( $v$ )
2:   while true do
3:     for every memory  $m$  in parallel do
4:        $Reg[m][i].T.write(ts)$ 
5:        $temp[m][1 \dots n] \leftarrow Reg[m][1 \dots n].read()$ 
6:     until completed for majority of memories
7:      $val \leftarrow temp[1..m][1..n].highestTspValue()$ 
8:     if  $val = \perp$  then  $val \leftarrow v$ 
9:     for every memory  $m$  in parallel do
10:       $Reg[m][i].V.write(val, ts)$ 
11:       $temp[m][1 \dots n] \leftarrow Reg[m][1..n].read()$ 
12:     until completed for majority of memories
13:     if  $ts = temp[1 \dots m][1 \dots n].highestTsp()$  then return ( $val$ )
14:      $ts \leftarrow ts + n$ 

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**Problem 2.** Consider the following variant of the Non-equivocating Broadcast algorithm seen today in class. Does this algorithm satisfy the Non-Equivocating Broadcast properties? Why or why not?

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**Algorithm 2** Non-equivocating Broadcast
 

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1: procedure BROADCAST( $m$ )
2:    $R[s].write(m)$ 
3: procedure RECEIVE
4:    $senderMsg = R[s].read()$ 
5:   for  $i = 1 \dots n$  do
6:      $recvMsg = R[i].read()$ 
7:     if  $recvMsg \neq \perp \wedge recvMsg \neq senderMsg$  then
8:        $\triangleright$  found conflicting values (Byzantine sender), don't deliver
9:     return
10:   $R[i].write(senderMsg)$ 
11:   $deliver(senderMsg)$ 

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