

Solution for Exercise 3

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LPD, EPFL

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The Splitter Object

- Only one operation: *splitter*
- Returns: *stop*, *left* or *right*
- If a **single** process executes *splitter*, then this process gets *stop*.
- If **two or more** processes invoke *splitter*, then not all get the same output.
- At most one process gets *stop*.

An Implementation of a Splitter

We use two registers:

- P (multi-valued), and
- S (binary, initialized to *false*)

upon *splitter* _{i}

```
 $P \leftarrow i;$ 
```

```
if  $S$  then return "right";
```

```
 $S \leftarrow true;$ 
```

```
if  $P = i$  then return "stop";
```

```
return "left";
```

An Implementation of a Splitter

We use two registers:

- P (multi-valued), and
- S (binary, initialized to *false*)

upon *splitter* _{i}

$P \leftarrow i$;

if S **then return** *"right"*;

$S \leftarrow true$;

if $P = i$ **then return** *"stop"*;

return *"left"*;

Non-adaptive Snapshot

upon *scan*;

$t_1 \leftarrow \text{collect}(), t_2 \leftarrow t_1;$

while *true* **do**

$t_3 \leftarrow \text{collect}();$

if $t_3 = t_2$ **then** **return** $\langle t_3[1].\text{val}, \dots, t_3[N].\text{val} \rangle ;$

for $k \leftarrow 1$ **to** N **do**

if $t_3[k].\text{ts} \geq t_1[k].\text{ts} + 2$ **then** **return** $t_3[k].\text{snapshot};$

$t_2 \leftarrow t_3;$

procedure *collect*()

for $k \leftarrow 1$ **to** N **do**

$x[k] \leftarrow R[k];$

return $x;$

Non-adaptive Snapshot (2)

```
procedure updatei(v)  
  ts ← ts + 1;  
  snapshot ← scan();  
  R[i] ← ⟨ ts, v, snapshot ⟩;
```

Adaptive Update

```
procedure update(v)  
  if myreg =  $\perp$  then  
    myreg  $\leftarrow$  obtain();  
  
  ts  $\leftarrow$  ts + 1;  
  snapshot  $\leftarrow$  scan();  
  R[myreg]  $\leftarrow$   $\langle$  ts, v, snapshot  $\rangle$ ;
```

Adaptive Scan

upon *scan*;

$t_1 \leftarrow \text{collect}(), t_2 \leftarrow t_1;$

while *true* **do**

$t_3 \leftarrow \text{collect}();$

if $t_3 = t_2$ **then return** $\langle t_3[1].\text{val}, \dots, t_3[t_3.\text{length}].\text{val} \rangle ;$

for $k \leftarrow 1$ **to** $t_3.\text{length}$ **do**

if $t_3[k].\text{ts} \geq t_1[k].\text{ts} + 2$ **then return** $t_3[k].\text{snapshot};$

$t_2 \leftarrow t_3;$

A Disallowed Solution

```
procedure obtain()  
   $\lfloor$  myreg  $\leftarrow$  C.fetch&inc();  
  
procedure collect()  
  for k  $\leftarrow$  1 to C.read() do  
     $\lfloor$  x[k]  $\leftarrow$  R[k];  
  return x;
```

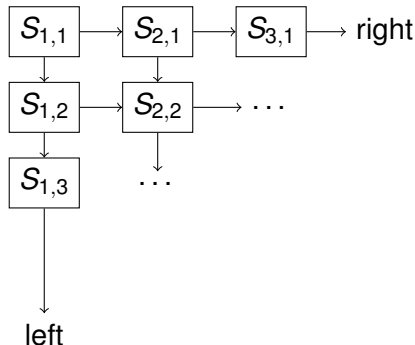
But we can use **only registers!**

The Splitter Object

- One operation: *splitter*
- Returns: *stop*, *left* or *right*
- If a **single** process executes *splitter*, then *stop* is returned.
- If **two or more** processes invoke *splitter*, then not all get the same output.
- At most one process gets *stop*.

Main Idea of Adaptive Snapshot

- Matrix of **registers** and **splitters**
- To obtain a register, a process must find a splitter that returns *stop*.
- Process starts from left top corner and follows the output of splitters.



The Obtain Operation

```
procedure obtain()  
   $x \leftarrow 1, y \leftarrow 1;$   
  while true do  
     $s \leftarrow S[x, y].\text{splitter}();$   
    if  $s = \text{"stop"}$  then  
       $\text{myreg} \leftarrow \langle x, y \rangle;$   
      return;  
    else if  $s = \text{"left"}$  then  $y \leftarrow y + 1;$   
    else  $x \leftarrow x + 1;$ 
```

The Collect Operation

procedure *collect*

$C \leftarrow \langle \rangle;$

$d \leftarrow 1;$

while *diagonal* d *has a splitter*
that has been traversed **do**

$C \leftarrow C \cdot \langle$ values of all
non- \perp registers on
diagonal d $\rangle;$

$d \leftarrow d + 1;$

return $C;$

